

DDH & attacks from genus theory

E/\mathbb{F}_q

$\text{end}(E) = \mathcal{O}$ order in $\mathcal{O}(\sqrt{t^2 - 4q})$

\mathbb{F}_q

$\#E(\mathbb{F}_q) = q + 1 - t$

$t^2 - 4q$ disc of $x^2 - tx + q$ ←

$\mathcal{O} \cong \mathbb{Z}[\pi]$ π Frobenius, root of

$\text{Ell}_q(\mathcal{O}, t) = \{ E' / \mathbb{F}_q : \text{end}(E') = \mathcal{O}$

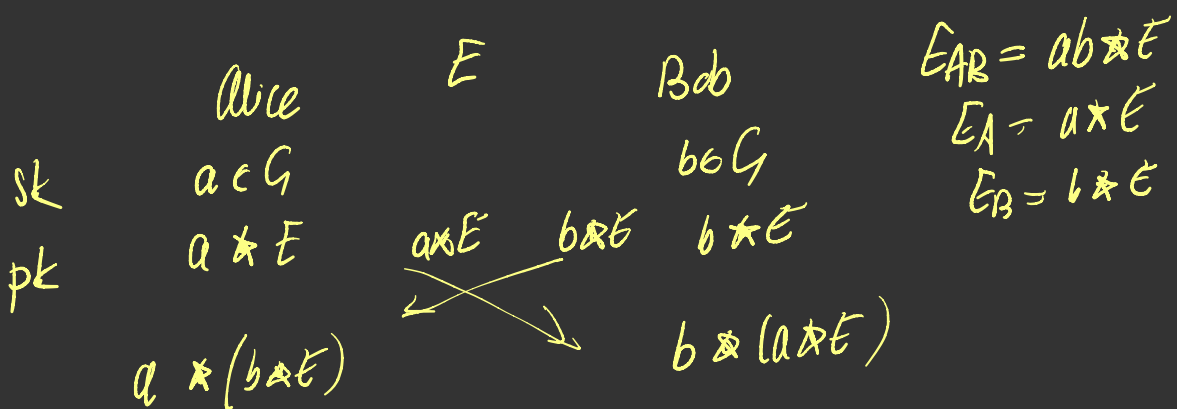
$\text{tr}(E') = t$

$\#E'(\mathbb{F}_q) = q + 1 - t$

$\left. \begin{matrix} \text{end}(E') = \mathcal{O} \\ \text{tr}(E') = t \\ \#E'(\mathbb{F}_q) = q + 1 - t \end{matrix} \right\} \mathbb{F}_q\text{-iso}$

$G = \mathcal{O}^\times$ acts on the set $X = \text{Ell}_q(\mathcal{O}, t)$

Group G acting on a set X , free and transitive



Secure if Alice and Bob are the only ones
who know E_{AB}

↳ CDH assumption

Problem $E, E_A, E_B \rightsquigarrow$ we can derive
information about E_{AB}

$$y^2 = x^3 + Ax^2 + x$$

$j(E_{AB})$ as shared key

DDH: there's nothing predictable about E_{AB}
if we only see E, E_A, E_B .

CSV: ^{sometimes} there exist characters χ st.

$$\chi(E_{AB}) = \chi(E_A) \cdot \chi(E_B)$$

χ character on $\mathcal{O}(\mathcal{D})$

χ quadratic character

$$\S 7 \quad \frac{\text{Pairings}}{E \times E} \quad E/\mathbb{F}_q \rightarrow \mathbb{F}_q$$

bilinear maps Weil pairing $E[\mathbb{N}] \times E[\mathbb{N}] \rightarrow \mu_{\mathbb{N}}$

$$e_n(P, Q) = \sum_{\mathbb{N}} e \in \mathbb{F}_q$$

m odd prime

$$E \rightarrow E' = [a] * E$$

$$[a] \in \mathcal{O}(\mathcal{O})$$

for every $a \in [a]$, we have an isogeny

$$\varphi_a: E \rightarrow E' \quad \deg \varphi_a = \text{norm}(a)$$

If we can say anything about $\deg \varphi_a \pmod m$

we get information about every $a \in [a]$

$$(\deg \varphi_a \pmod m)_{a \in [a]}$$

E, E' isogenous curves $\text{val}_m(\#E(\mathbb{F}_2)) = 1$

$$P \in E(\mathbb{F}_2)[m]$$

$$P' \in E'(\mathbb{F}_2)[m]$$

then degree of any isogeny

$$\varphi: E \rightarrow E'$$

is, up to squares mod m

$$\deg \varphi = \left(\log_{T_m(P, \rho)} \log |T_m(P', P')| \right)$$

$$\begin{array}{ccc} E & \xrightarrow{a} & E_A \\ b \downarrow & & \downarrow b \\ E_B & \xrightarrow{a} & E_{AB} \cong E_C \end{array}$$

$G = d(\mathcal{O})$ \mathcal{O} order in im. quad field

$X = \text{ell}$ \leftarrow elliptic curves / \mathbb{F}_q

$$\text{End}_{\mathbb{F}_q}(E) \cong \mathcal{O}$$

trace $t = \#E(\mathbb{F}_q) - q - 1$ is fixed

$G \curvearrowright X$ is free and transitive

DH key exchange

E = starting pt

Alice

Bob

sk

a

b

pk

$a * E$

$b * E$



$a * (b * E)$

$b * (a * E)$

1) Secure if Alice and Bob the only ones } CDH
who hold $E_{AB} = ab * E$

DDA $\approx E, E_A, E_B$ then E_{AB} "looks random"
 $\swarrow \searrow$
 public keys

you can't use E, E_A, E_B to predict anything
 about E_{AB}

How many supersingular curves are there

$\approx \frac{p}{12} = \# \text{supersingular } j\text{-invariants in } \mathbb{F}_p$

$\sqrt{p} = \# \text{supersingular curves over } \mathbb{F}_p, \text{ up to iso}$

$\approx \# \text{cl}(\mathbb{Z}[\sqrt{-p}])$

$$h(-\Delta) \approx \sqrt{|\Delta|}$$

Winter week 3

$p \bmod 8$	
1 mod 4	$h(-4p)$
3 mod 8	$h(-4p) + h(-p)$
5 mod 8	$2 \cdot h(-p)$

Ex 1.1. there are p different A 's
 only \sqrt{p} of them are supersingular curves $\left\{ \frac{\sqrt{p}}{p} \right\}$

§2 Pairings m odd prime, E/\mathbb{F}_q

$$T_m: E(\mathbb{F}_q)[m] \times E(\mathbb{F}_q) \xrightarrow{\mu_m} \mu_m$$

$\mu_m \subseteq \mathbb{F}_q$
 $\Leftrightarrow m|q-1$

$$(P, Q) \mapsto T_m(P, Q)$$

$$\varphi: E \rightarrow E'$$
$$P, Q \mapsto \varphi(P), \varphi(Q)$$

$$T_m(\varphi(P), \varphi(Q)) = T_m(P, Q)^{\deg \varphi}$$

$$\mapsto \deg \varphi \pmod{m}$$

"Hard problem": $E \rightarrow E'$ under some secret isogeny φ

find pair $P, P' = \varphi(P)$

Get-around: isogenies

$\left\{ \begin{array}{l} \text{rational map: } \text{points in } E(\mathbb{F}_q) \text{ map to } E'(\mathbb{F}_q) \\ \text{group homo: } E[m] \longrightarrow E'[m] \end{array} \right.$

$$\begin{array}{l} 1) \quad \varphi: E(\mathbb{F}_q)[m] = \langle P \rangle \\ \quad \quad E'(\mathbb{F}_q)[m] = \langle P' \rangle \end{array} \implies \varphi(P) = \ell \cdot P' \quad \ell \text{ unknown}$$

$E \qquad \qquad E'$

(conditions apply)

computing these Tate pairings determines

$\deg \varphi \pmod m$, up to squares
 $\pmod m$

for any isogeny $\varphi: E \rightarrow E'$

DDH If all these conditions hold, and we can compute

$$E \xrightarrow{\alpha} E_A$$

deg of isog $(E \rightarrow E_A)$ mod m up to squares

$$\begin{array}{ccc} \downarrow & & \downarrow \\ E_B & \xrightarrow{\alpha} & E_{AB} \stackrel{?}{=} E_C \quad E_C \text{ random} \end{array}$$

deg of isog $(E_B \rightarrow E_{AB}) \dots$

$$\chi(E, E_A) = \left(\frac{\deg \varphi}{m} \right) \quad \leftarrow \text{Legendre symbol}$$

$$= \begin{cases} 1 & \text{if } \deg \varphi \text{ always square mod } m \\ -1 & \text{if } \deg \varphi \text{ always non-square} \end{cases}$$

always: for any isogeny $\varphi: E \rightarrow E'$

Conditions: $\mu_m \subseteq \mathbb{F}_q$, $E(\mathbb{F}_q)[m^\infty] = \langle P \rangle$ P of order m

SIDH

$$E \xrightarrow{2^n} E_A$$

$$3^m \downarrow \\ E_B$$

$$E_C \stackrel{?}{=} E_{AB}$$

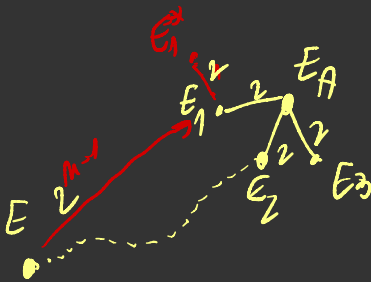
E_{AB} is 3^m -isogenous to E_A
 2^n -isogenous to E_B

- A) Decide whether E_C is 3^m -isog to E_A
- B) 2^n -isog to E_B

CSIDH
 $E \cong E'$
 always
 $E' = a \times E$
 for some a

- I) Galbraith - Vercauteren Comp problems in log...
- II) Jao - Urbanik? Sok

DDH to CDH reduction isogeny finding



$m = \# \text{ steps in SIDH}$
 $= \text{poly size}$
 $= O(\log p)$

Input to these problems

$$\mathbb{P}, \# \mathbb{P}^2, E_0 / \mathbb{F}_p^2, \dots$$

$$2^n \approx \sqrt{p}$$

$$m \approx \frac{1}{2} \log p$$

write it down using $\log_2 p$ bits

Polynomial time = poly in $\log_2 p$